



Xpress Performance Workshop RCRblackboard

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RESEARCH \ ENGAGEMENT \ INNOVATION

Why RCRToolkit?

- Thread performance on multi-core systems limited by what the other threads on the system are simultaneously doing.
 - L3 cache contention
 - Memory bandwidth limitations – includes contention in DIMMs
 - Internal bus contention
- RENCI has been working on Resource Centric Reflection toolkit to expose contention to programmer and runtime.
 - Performance Tuning Tools (RCRoolkit)
 - Runtime Scheduler (MAESTRO scheduler in Qthreads)

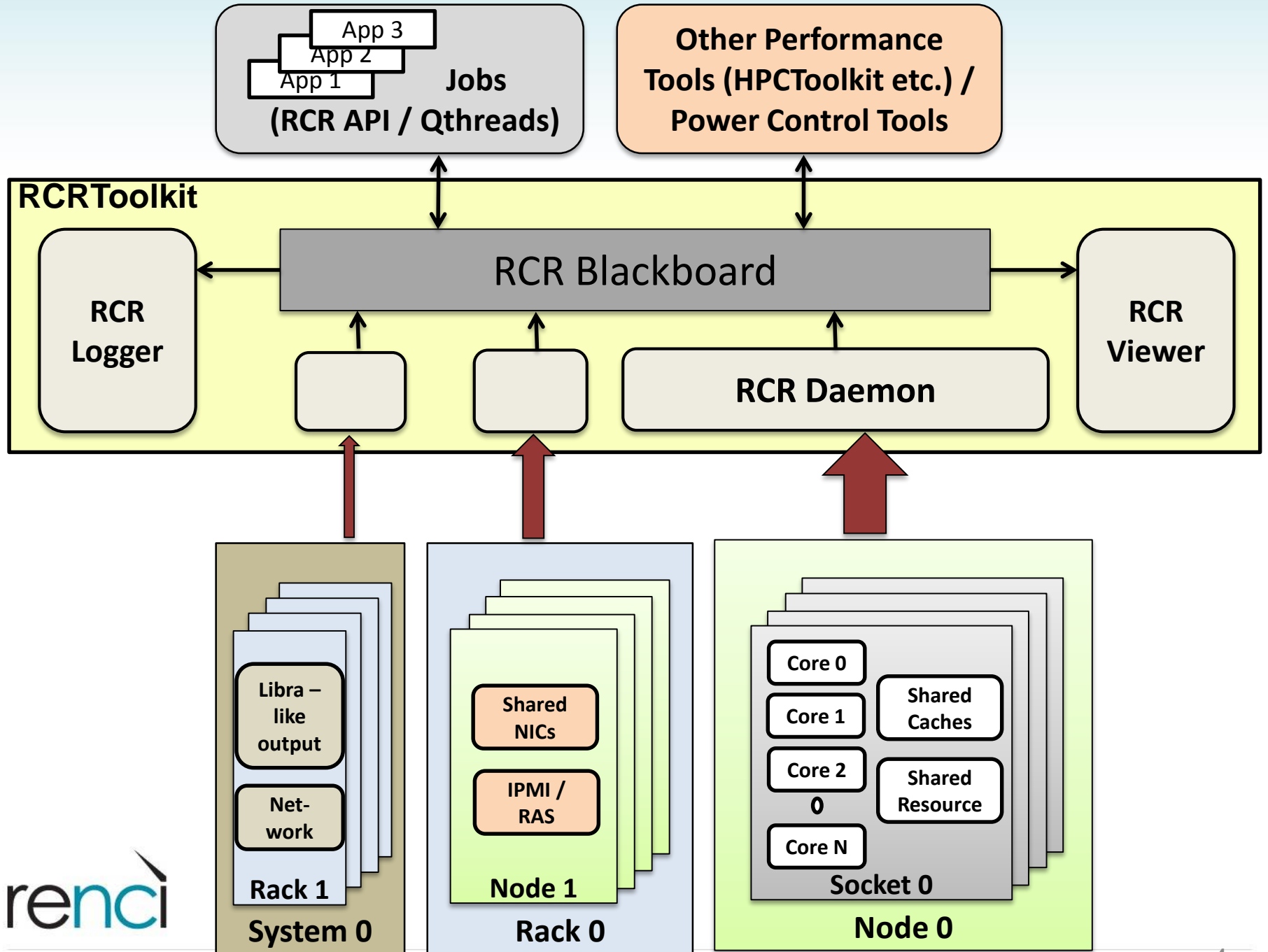
Pieces of RCRToolkit

- Infrastructure

- RCRblackboard – database to store dynamic information about system
- RCRdaemon - allow user access to hardware performance counters
- RCRlogger – allow post-execution review of counters
- RCRviewer – simple GUI to view results

- Clients

- EnergyStat API – allow user to see energy application required
- Qthread scheduler – allow runtime access to dynamic information
- HPCToolkit – Modified to query blackboard



RCRblackboard (1)

- **Publisher/Reader Semantics**
 - Each section 1 writer multiple readers – eliminate synchronization
 - No reader checkin – writer does not produce events for readers
 - Self-describing data format that writer/readers agree on
- **Uses shared memory regions**
 - One per writer
 - currently only one writer – it uses /dev/shm/bbFile

RCRblackboard (2)

- Google Protobuf

- Self-describing, compact
- Seems designed for network and stores in a compressed format
 - Compression on every write is very expensive for us
 - Future – write store function that doesn't compress
 - Updates become simple writes / no compression
 - Reads are simple reads / no expansion
 - Will need to define mechanism to prevent inconsistent data being read (when reading multiple values – 2 version numbers?)
- Hierarchical based on classes from protoc
 - On our 2 socket SandyBridge system
 - 8 sets of core counters
 - 2 sets of socket counters
 - System-wide counters

RCRblackboard (3) -- partial protoc def

from protobuf/blackboard.proto

```
Message RCRBlackboard {  
    optional RCRBlackboardMetadata bbMetadata = 1  
    repeated RCRNode node = 2  
    repeated RCRSocket socket = 3  
    repeated RCRCore core = 4  
    repeated RCRSocketMeter socketMeter = 5  
    repeated RCRCoreMeter coreMeter = 6  
}
```

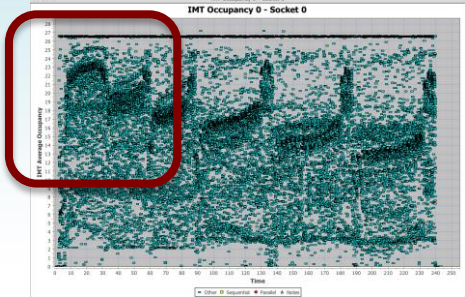

RCRdaemon

- Write hardware counters into RCRblackboard
 - Chip-wide – energy/L3 cache/Memory Controller
 - Core-specific – std set (cycle cnt/floating pt/etc.)
- Several Architecture specific versions
 - Intel SandyBridge (currently used)
 - Intel Nahalem (compiles – as of Monday / untested – doesn't crash immediately)
 - AMD Opteron (used in the past and probably victim of bit rot)
- Needs to run at kernel protection level to access global counters
 - Configuration dependent (energy counter requires it / as do some L3 counters)
- Writes `/dev/shm/bbFile` using protobuf interface
 - Current overhead ~16% of one core
 - Big savings by eliminating compression (one per write)

RCRlogger

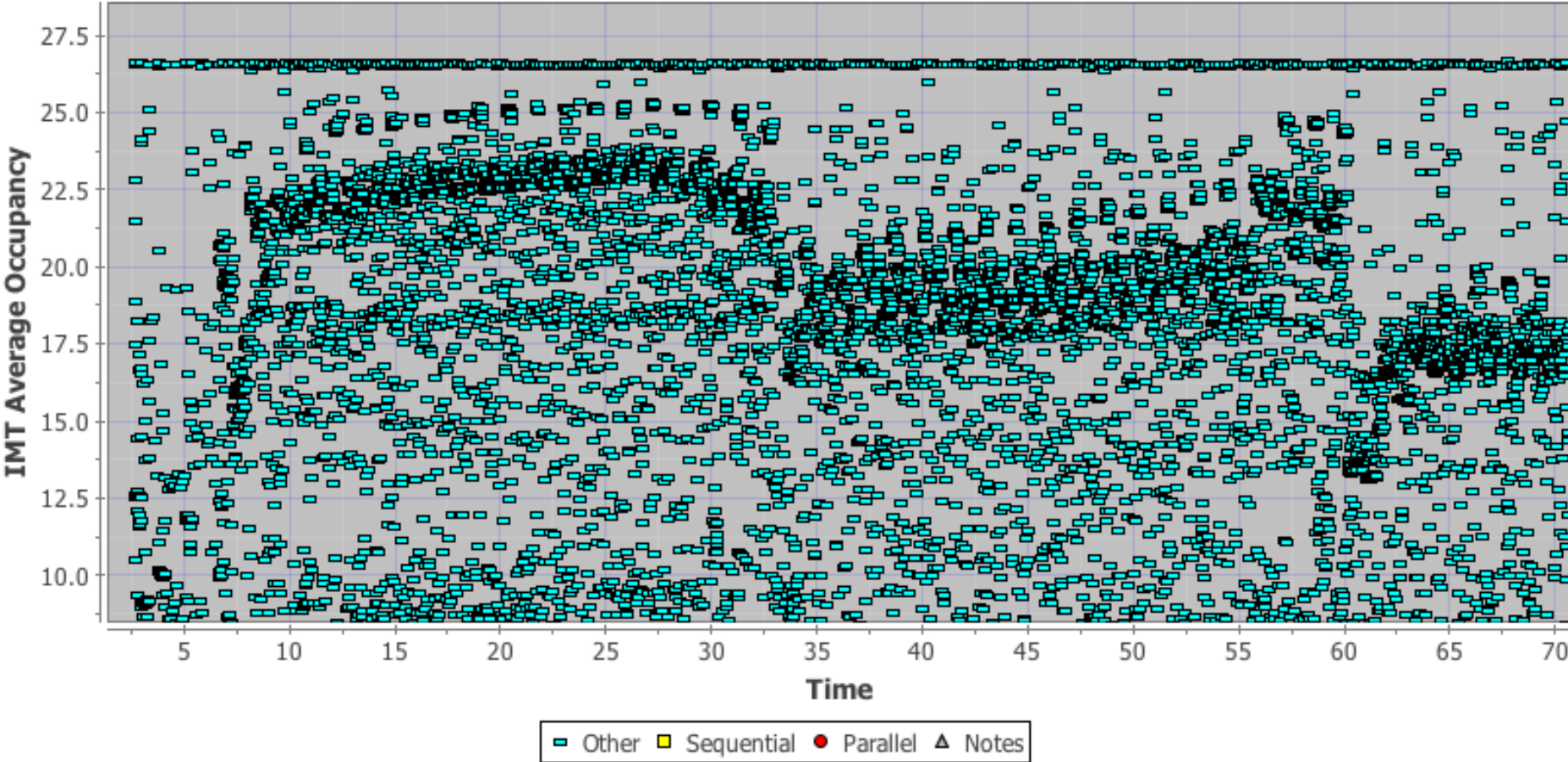
- Reads RCRblackboard periodically and writes results to stdout
 - Dumps all active counters on single line
 - Identified by socket/core number and counter number
 - Up to ~12000 times a sec on Intel SandyBridge (2.7GHz)
 - Faster than many of the counters update in RCRblackboard (energy ~1000)
 - Startup option to set frequency (-i #in microseconds)
 - -d turns into daemon (no stdout – not sure why)
 - No -f output to filename (should be added)

RCRviewer



IMT Occupancy 0 - Socket 0

IMT Occupancy 0 - Socket 0



EnergyStat API

- Provides a pair of calls (in C) to capture energy usage during a program
 - extern “C” int energyDaemonInit(int wait);
 - extern “C” void energyDaemonTerm();
- Produces these lines of output
 - (init call) Starting doEnergyWork
 - (term call) Application (Energy) – Time 8.109619 Total energy consumed 1072.728810 Ave. Power Level 132.278572 Final Temperature socket 1 – 53.000000 socket 2 – 46.000000
- Multiple calls to energyDaemonTerm allowed
 - Each prints energy since previous call
- Initiates low-overhead daemon
 - Wakes up every wait nanoseconds and reads counters (32 bit – protects from overflow)
 - Only works on Intel SandyBridge (and probably IvyBridge) processor

Sherwood Scheduler

- Locality-aware scheduler for Qthreads
 - Work sharing between cores sharing L3 cache
 - Work stealing between sockets sharing an address space
- Modified to reduce energy consumption
 - Reads energy and memory concurrency from RCRblackboard
 - If both high reduces the number of active threads
 - Duty-cycle modification to greatly reduce power requirements of idle threads
 - Saved ~3% power for benchmarks/Mini-Apps where it applied
 - Micro-algorithm benchmarks(UNC), BOTS suite(Barcelona), and LULESH (LLNL mini-app)

HPCToolkit hot-wired with RCRToolkit

hpcviewer: fft_open_mp

```

483  _p_jc = (p_index_ * mj2);
484  _p_jd = _p_jc;
485  _p_wjw[0] = ( *w)[[_p_jw * 2] + 0];
486  _p_wjw[1] = ( *w)[[_p_jw * 2] + 1];
487  if ( *sgn < 0.0) {
488  _p_wjw[1] = -_p_wjw[1];
489  }
490  for (_p_k = 0; _p_k < *mj; _p_k++) {
491  ( *c)[[_p_jc + _p_k * 2] + 0] = (( *a)[[_p_ja + _p_k * 2] + 0] + ( *b)[[_p_jb + _p_k * 2] + 0]);
492  ( *c)[[_p_jc + _p_k * 2] + 1] = (( *a)[[_p_ja + _p_k * 2] + 1] + ( *b)[[_p_jb + _p_k * 2] + 1]);
493  _p_ambr = (( *a)[[_p_ja + _p_k * 2] + 0] - ( *b)[[_p_jb + _p_k * 2] + 0]);
494  _p_ambu = (( *a)[[_p_ja + _p_k * 2] + 1] - ( *b)[[_p_jb + _p_k * 2] + 1]);
495  ( *d)[[_p_jd + _p_k * 2] + 0] = (_p_wjw[0] * _p_ambr) - (_p_wjw[1] * _p_ambu);
496  ( *d)[[_p_jd + _p_k * 2] + 1] = (_p_wjw[1] * _p_ambr) + (_p_wjw[0] * _p_ambu);
497  }
498  }
499  }
500
501 stat
502 {
503 do
504 do
505 int *n2 = (int *)(((struct OUT__2_1527__data *)__out_argv) -> OUT__2_1527__data::n2_p);

```

PAPI_TOT_CYC:Sum (l)	PAPI_L2_TCM:Sum (l)	RCR-PAPI_L2_TCM:Sum (l)	Contention Percentage
1.37e+12 100 %	2.69e+09 100 %	1.46e+09 100 %	5.43e+01

Calling Context View | Callers View | Flat View

Scope	PAPI_TOT_CYC:Sum (l)	PAPI_L2_TCM:Sum (l)	RCR-PAPI_L2_TCM:Sum (l)	Contention Percentage
Experiment Aggregate Metrics	1.37e+12 100 %	2.69e+09 100 %	1.46e+09 100 %	5.43e+01
monitor main	8.40e+11 61.2%	1.26e+09 46.8%	3.50e+08 24.0%	2.78e+01
atthread wrapper	4.91e+11 35.8%	1.25e+09 46.5%	1.06e+09 72.6%	8.48e+01
qloop step wrapper	4.90e+11 35.7%	1.23e+09 45.7%	1.04e+09 71.2%	8.46e+01
OUT_1_1527	4.71e+11 34.3%	1.22e+09 45.4%	1.04e+09 71.2%	8.52e+01
loop at fft_open_mp.rose: 479	4.66e+11 33.9%	1.22e+09 45.4%	1.04e+09 71.2%	8.52e+01
loop at fft open mp.rose: 490	4.57e+11 33.3%	1.20e+09 44.6%	1.02e+09 69.9%	8.50e+01
fft open mp.rose: 486	3.77e+09 0.3%			
inlined from fft open mp.rose: 454	3.15e+09 0.2%	2.00e+07 0.7%	2.00e+07 1.4%	1.00e+02
fft open mp.rose: 487	7.74e+08 0.1%			
fft open mp.rose: 479	3.58e+08 0.0%			
fft open mp.rose: 485	3.13e+08 0.0%			
fft open mp.rose: 488	3.60e+07 0.0%			
XOMP loop default	5.37e+09 0.4%			
loop at fft open mp.rose: 490	4.57e+11 33.3%	1.20e+09 44.6%	1.02e+09 69.9%	8.50e+01
OUT_3_1527	4.83e+09 0.4%			

321M of 513M